

Low Rank Approximation

Lecture 6

Daniel Kressner

Chair for Numerical Algorithms and HPC

Institute of Mathematics, EPFL

`daniel.kressner@epfl.ch`



ÉCOLE POLYTECHNIQUE
FÉDÉRALE DE LAUSANNE



Tensor Train (TT) decomposition

A tensor \mathcal{X} is in **TT decomposition** if it can be written as

$$\mathcal{X}(i_1, \dots, i_d) = \sum_{k_1=1}^{r_1} \cdots \sum_{k_{d-1}=1}^{r_{d-1}} \mathcal{U}_1(1, i_1, k_1) \mathcal{U}_2(k_1, i_2, k_2) \cdots \mathcal{U}_d(k_{d-1}, i_d, 1).$$

- ▶ Smallest possible tuple (r_1, \dots, r_{d-1}) is called **TT rank** of \mathcal{X} .
- ▶ $\mathcal{U}_\mu \in \mathbb{R}^{r_{\mu-1} \times n_\mu \times r_\mu}$ (formally set $r_0 = r_d = 1$) are called **TT cores** for $\mu = 1, \dots, d$.
- ▶ If TT ranks are not large \rightsquigarrow high compression ratio as d grows.
- ▶ TT decomposition multilinear wrt cores.
- ▶ TT decomposition connects to
 - ▶ matrix products \rightsquigarrow **Matrix Product States** (MPS) in physics (see [Grasedyck/Kressner/Tobler'2013] for references)
 - ▶ simultaneous matrix factorizations \rightsquigarrow SVD-based compression
 - ▶ contractions and tensor network diagrams \rightsquigarrow design of efficient contraction-based algorithms

TT decomposition and matrix products

$$\mathcal{X}(i_1, \dots, i_d) = \sum_{k_1=1}^{r_1} \cdots \sum_{k_{d-1}=1}^{r_{d-1}} \mathcal{U}_1(1, i_1, k_1) \mathcal{U}_2(k_1, i_2, k_2) \cdots \mathcal{U}_d(k_{d-1}, i_d, 1).$$

Let $U_\mu(i_\mu)$ be i_μ th slice of μ th core: $U_\mu(i_\mu) := \mathcal{U}_\mu(:, i_\mu, :) \in \mathbb{R}^{r_{\mu-1} \times r_\mu}$.
Then

$$\mathcal{X}(i_1, i_2, \dots, i_d) = U_1(i_1) U_2(i_2) \cdots U_d(i_d).$$

EFY. Given two tensors \mathcal{X} and \mathcal{Y} in TT decomposition, derive a TT decomposition for $\mathcal{X} + \mathcal{Y}$.

EFY. Show that a tensor of tensor rank R has TT rank at most (R, \dots, R) by converting its CP decomposition into a TT decomposition.

EFY. Given a tensor in TT decomposition, derive upper bounds for its multilinear rank.

Remark: Error analysis of matrix multiplication [Higham'2002] shows that TT decomposition may suffer from numerical instabilities if

$$\|U_1(i_1)\|_2 \|U_2(i_2)\|_2 \cdots \|U_d(i_d)\|_2 \gg |\mathcal{X}(i_1, i_2, \dots, i_d)|.$$

See [Bachmayr/Kazeev: arXiv:1802.09062] for more details.

TT decomposition and matrix factorizations

$$\mathcal{X}(i_1, \dots, i_d) = \sum_{k_1, k_2, \dots, k_{d-1}} \mathcal{U}_1(1, i_1, k_1) \mathcal{U}_2(k_1, i_2, k_2) \cdots \mathcal{U}_d(k_{d-1}, i_d, 1).$$

For any $1 \leq \mu \leq d - 1$ group first μ factors and last $d - \mu$ factors together:

$$\begin{aligned} & \mathcal{X}(i_1, \dots, i_\mu, i_{\mu+1}, \dots, i_d) \\ = & \sum_{k_\mu=1}^{r_\mu} \left(\sum_{k_1, \dots, k_{\mu-1}} \mathcal{U}_1(1, i_1, k_1) \cdots \mathcal{U}_\mu(k_{\mu-1}, i_\mu, k_\mu) \right) \\ & \cdot \left(\sum_{k_{\mu+1}, \dots, k_{d-1}} \mathcal{U}_{\mu+1}(k_\mu, i_{\mu+1}, k_{\mu+1}) \cdots \mathcal{U}_d(k_{d-1}, i_d, 1) \right) \end{aligned}$$

This can be interpreted as a matrix-matrix product of two (large) matrices!

TT decomposition and matrix factorizations

The μ th unfolding of $\mathcal{X} \in \mathbb{R}^{n_1 \times n_2 \times \dots \times n_d}$ is obtained by arranging the entries in a matrix

$$\mathcal{X}^{<\mu>} \in \mathbb{R}^{(n_1 n_2 \dots n_\mu) \times (n_{\mu+1} \dots n_d)}$$

where the corresponding index map is given by

$$\iota : \mathbb{R}^{n_1 \times \dots \times n_d} \rightarrow \mathbb{R}^{n_1 \dots n_\mu} \times \mathbb{R}^{n_{\mu+1} \dots n_d}, \quad \iota(i_1, \dots, i_d) = (i_{\text{row}}, i_{\text{col}}),$$

$$i_{\text{row}} = 1 + \sum_{\nu=1}^{\mu} (i_\nu - 1) \prod_{\tau=1}^{\nu-1} n_\tau, \quad i_{\text{col}} = 1 + \sum_{\nu=\mu+1}^d (i_\nu - 1) \prod_{\tau=\mu+1}^{\nu-1} n_\tau.$$

TT decomposition and matrix factorizations

Define **interface matrices**

$$\mathbf{X}_{\leq \mu} \in \mathbb{R}^{n_1 n_2 \cdots n_\mu \times r_\mu}, \quad \mathbf{X}_{\geq \mu+1} \in \mathbb{R}^{n_{\mu+1} n_{\mu+2} \cdots n_d \times r_\mu}$$

as

$$\mathbf{X}_{\leq \mu}(i_{\text{row}}, j) = \sum_{k_1, \dots, k_{\mu-1}} \mathcal{U}_1(1, i_1, k_1) \cdots \mathcal{U}_{\mu-1}(k_{\mu-2}, i_{\mu-1}, k_{\mu-1}) \mathcal{U}_\mu(k_{\mu-1}, i_\mu, j)$$

$$\mathbf{X}_{\geq \mu+1}(i_{\text{col}}, j) = \sum_{k_{\mu+1}, \dots, k_{d-1}} \mathcal{U}_{\mu+1}(j, i_{\mu+1}, k_{\mu+1}) \mathcal{U}_{\mu+2}(k_{\mu+1}, i_{\mu+2}, k_{\mu+2}) \cdots \mathcal{U}_d(k_{d-1}, i_d, 1)$$

Matrix factorizations

$$\mathbf{X}^{<\mu>} = \mathbf{X}_{\leq \mu} \mathbf{X}_{\geq \mu+1}^T, \quad \mu = 1, \dots, d-1.$$

TT decomposition and matrix factorizations

Important: These matrix factorizations are nested!

$$X_{\leq \mu} = (I_{n_\mu} \otimes X_{\leq \mu-1}) U_\mu^L, \quad \text{and} \quad X_{\geq \mu}^T = U_\mu^R (X_{\geq \mu+1}^T \otimes I_{n_\mu}),$$

where

$$U_\mu^L = U_\mu^{<2>}, \quad U_\mu^R = U_\mu^{(1)} = U_\mu^{<1>}.$$

The relations $X_{\leq 1} = U_1 \equiv \mathcal{U}_1$ and

$$X_{\leq \mu} = (I_{n_\mu} \otimes X_{\leq \mu-1}) U_\mu^L, \quad \mu = 2, \dots, d,$$

fully characterize the TT decomposition:

$$\begin{aligned} \text{vec}(\mathcal{X}) &= X_{\leq d} \\ &= (I \otimes X_{\leq d-1}) U_d^L = (I \otimes X_{\leq d-1}) \text{vec}(\mathcal{U}_d) \\ &= (I \otimes I \otimes X_{\leq d-2}) (I \otimes U_{d-1}^L) \text{vec}(\mathcal{U}_d) \\ &\vdots \\ &= (I \otimes \dots \otimes I \otimes U_1^L) \dots (I \otimes U_{d-1}^L) \text{vec}(\mathcal{U}_d) \end{aligned}$$

EFY. Perform an analogous calculation for $X_{\geq \mu}^T$, that is, resolve the recursion $X_{\geq \mu}^T = U_\mu^R (X_{\geq \mu+1}^T \otimes I_{n_\mu})$.

TT decomposition and matrix factorizations

Lemma

The TT rank of a tensor is given by

$$(\text{rank } X^{<1>}, \dots, \text{rank } X^{<d-1>})$$

Proof. Because of the connection to matrix factorizations, the TT rank cannot be smaller than $(\text{rank } X^{<1>}, \dots, \text{rank } X^{<d-1>})$. We need to exclude that it can be larger. For this purpose, we construct a TT decomposition with

$$(r_1, \dots, r_{d-1}) := (\text{rank } X^{<1>}, \dots, \text{rank } X^{<d-1>}).$$

Step 1: Factorize

$$X^{<1>} = U_1 \tilde{X}^{<1>}, \quad U_1 \in \mathbb{R}^{n_1 \times r_1}, \quad \tilde{X}^{<1>} \in \mathbb{R}^{r_1 \times n_2 \cdots n_d},$$

and hence

$$\tilde{X}^{<1>} = U_1^\dagger X^{<1>}, \quad U_1^\dagger = (U_1^T U_1)^{-1} U_1^T$$

In terms of tensors: $\mathcal{X} = U_1 \circ_1 \tilde{\mathcal{X}}$.

$U_1 \equiv U_1$ is the first TT core (and $X_{\geq 1}^T := \tilde{X}^{<1>}$).

Relation for second unfolding via Kronecker product:

$$X^{<2>} = (I_{n_2} \otimes U_1) \tilde{X}^{<2>}.$$

Together with full column rank of U_1 , this implies

$$\text{rank}(\tilde{X}^{<2>}) = \text{rank}(X^{<2>}) = r_2.$$

Step 2: Factorize

$$\tilde{X}^{<2>} = U_2^L \hat{X}^{<1>}, \quad U_2^L \in \mathbb{R}^{r_1 n_2 \times r_2}, \quad \hat{X}^{<1>} \in \mathbb{R}^{r_2 \times n_3 \cdots n_d},$$

U_2^L gives the second TT core $U_2 \in \mathbb{R}^{r_1 \times n_2 \times r_2}$ and $X_{\geq 2}^T := \hat{X}^{<1>}$.

Relation for third unfolding via Kronecker product:

$$X^{<3>} = (I_{n_3} \otimes I_{n_2} \otimes U_1) \tilde{X}^{<3>} = (I_{n_3} \otimes I_{n_2} \otimes U_1) (I_{n_3} \otimes U_2^L) \hat{X}^{<2>}$$

Together with full column ranks of U_1 and U_2^L , this implies

$$\text{rank}(\hat{X}^{<2>}) = \text{rank}(X^{<3>}) = r_3.$$

Continuing in this manner gives cores $U_\mu \in \mathbb{R}^{r_{\mu-1} \times n_\mu \times r_\mu}$ such that

$$\text{vec}(\mathcal{X}) = (I \otimes \cdots \otimes I \otimes U_1) \cdots (I \otimes U_{d-1}^L) \text{vec}(U_d)$$

This defines a TT decomposition.

Truncation in TT format

Proof of Lemma can be turned into practical algorithm (TT-SVD by [Oseledets'2011]) for approximating a given tensor \mathcal{X} in TT format:

Input: $\mathcal{X} \in \mathbb{R}^{n_1 \times \dots \times n_d}$, target TT rank (r_1, \dots, r_{d-1}) .

Output: TT cores $\mathcal{U}_\mu \in \mathbb{R}^{r_{\mu-1} \times n_\mu \times r_\mu}$ that define a TT decomposition approximating \mathcal{X} .

- 1: Set $r_0 = r_d = 1$. (and formally add leading singleton dimension to $\mathcal{X} \in \mathbb{R}^{1 \times n_1 \times \dots \times n_d}$).
- 2: **for** $\mu = 1, \dots, d - 1$ **do**
- 3: Reshape \mathcal{X} into $X^{<2>} \in \mathbb{R}^{r_{\mu-1} n_\mu \times n_{\mu+1} \dots n_d}$.
- 4: Compute rank- r_μ approximation $X^{<2>} \approx U \Sigma V^T$ (e.g., via SVD)
- 5: Reshape U into $\mathcal{U}_\mu \in \mathbb{R}^{r_{\mu-1} \times n_\mu \times r_\mu}$.
- 6: Update \mathcal{X} via $X^{<2>} \leftarrow U^T X^{<2>} = \Sigma V^T$.
- 7: **end for**
- 8: Set $\mathcal{U}_d = \mathcal{X}$.

Truncation in TT format

Theorem

Let \mathcal{X}_{SVD} denote the tensor in TT decomposition obtained from TT-SVD. Then

$$\|\mathcal{X} - \mathcal{X}_{\text{SVD}}\| \leq \sqrt{\varepsilon_1^2 + \dots + \varepsilon_d^2},$$

where

$$\varepsilon_\mu^2 = \|\mathcal{X}^{<\mu>} - \mathcal{T}_{r_\mu}(\mathcal{X}^{<\mu>})\|_F^2 = \sigma_{r_\mu+1}(\mathcal{X}^{<\mu>})^2 + \dots.$$

Proof. After μ steps of the algorithm we have the following situation:

- ▶ Core tensors $\mathcal{U}_1, \dots, \mathcal{U}_\mu$ have been computed, defining $\mathcal{X}_{\leq \mu}$.
- ▶ Remaining tensor has size $r_\mu \times n_{\mu+1} \times \dots \times n_d$.

Reshape remaining tensor into matrix $\mathbf{Y}_{\geq \mu}^T \in \mathbb{R}^{r_\mu \times n_{\mu+1} \dots n_d}$. Will prove relations

$$\mathbf{X}_{\leq \mu}^T \mathbf{X}_{\leq \mu} = \mathbf{I}, \quad \mathbf{Y}_{\geq \mu+1}^T = \mathbf{X}_{\leq \mu}^T \mathcal{X}^{<\mu>},$$

and

$$\|\mathcal{X}^{<\mu>} - \mathbf{X}_{\leq \mu} \mathbf{Y}_{\geq \mu+1}^T\|_F \leq \sqrt{\varepsilon_1^2 + \dots + \varepsilon_\mu^2} \quad (1)$$

for $\mu = 1, \dots, d-1$ by induction. For $\mu = d-1$, this shows the theorem.

Line 3 in the μ th step of the algorithm proceeds by reshaping the remaining tensor from step $\mu - 1$ (corresponding to $Y_{\geq \mu}^T$) into an array of size $Y^{<2>} \in \mathbb{R}^{r_{\mu-1} n_{\mu} \times n_{\mu+1} \cdots n_d}$. By induction assumption,

$$Y_{\geq \mu}^T = X_{\leq \mu-1}^T X^{<\mu-1>} \Rightarrow Y^{<2>} = (I_{n_{\mu}} \otimes X_{\leq \mu-1})^T X^{<\mu>}. \quad (2)$$

The matrix $U_{\mu}^L \equiv U$ computed in Line 4 contains left singular vectors of $Y^{<2>}$. In particular, $(U_{\mu}^L)^T U_{\mu}^L = I$. Together with the induction assumption and the relation

$$X_{\leq \mu} = (I_{n_{\mu}} \otimes X_{\leq \mu-1}) U_{\mu}^L,$$

this implies $X_{\leq \mu}^T X_{\leq \mu} = I$ and $Y_{\geq \mu+1}^T = X_{\leq \mu}^T X^{<\mu>}$. Moreover,

$$\begin{aligned} \|(I - U_{\mu}^L (U_{\mu}^L)^T) Y^{<2>}\|_F &= \|Y^{<2>} - \mathcal{T}_{r_{\mu}}(Y^{<2>})\|_F \\ &\leq \|X^{<\mu>} - \mathcal{T}_{r_{\mu}}(X^{<\mu>})\|_F = \varepsilon_{\mu}. \end{aligned}$$

Finally, we obtain:

$$\begin{aligned}
& \|X^{<\mu>} - X_{\leq\mu} X_{\leq\mu}^T X^{<\mu>} \|_F^2 \\
= & \|X^{<\mu>} - (I \otimes X_{\leq\mu-1}) U_\mu^L (U_\mu^L)^T (I \otimes X_{\leq\mu-1})^T X^{<\mu>} \|_F^2 \\
= & \|X^{<\mu>} - (I \otimes X_{\leq\mu-1}) (I \otimes X_{\leq\mu-1})^T X^{<\mu>} \|_F^2 \\
+ & \|(I \otimes X_{\leq\mu-1}) (I \otimes X_{\leq\mu-1})^T X^{<\mu>} \\
& - (I \otimes X_{\leq\mu-1}) U_\mu^L (U_\mu^L)^T (I \otimes X_{\leq\mu-1})^T X^{<\mu>} \|_F^2 \\
= & \|X^{<\mu-1>} - X_{\leq\mu-1} X_{\leq\mu-1}^T X^{<\mu-1>} \|_F^2 + \|(I - U_\mu^L (U_\mu^L)^T) Y^{<2>} \|_F^2 \\
\leq & \varepsilon_1^2 + \cdots + \varepsilon_{\mu-1}^2 + \varepsilon_\mu^2
\end{aligned}$$

This completes the proof.

Truncation in TT format

Consequence of the theorem:

Corollary

Let $\mathcal{X}_{\text{best}}$ denote the best approximation of \mathcal{X} with TT rank (r_1, \dots, r_{d-1}) . Then

$$\|\mathcal{X} - \mathcal{X}_{\text{SVD}}\| \leq \sqrt{d-1} \|\mathcal{X} - \mathcal{X}_{\text{best}}\|.$$

Proof.

Since $\mathcal{X}_{\text{best}}^{<\mu>}$ has rank r_μ we have

$$\begin{aligned} \|\mathcal{X} - \mathcal{X}_{\text{best}}\| &\geq \max_{\mu=1, \dots, r-1} \|\mathcal{X}^{<\mu>} - \mathcal{X}_{\text{best}}^{<\mu>}\|_F \\ &= \max_{\mu=1, \dots, r-1} \varepsilon_\mu \geq \frac{1}{\sqrt{d-1}} \sqrt{\varepsilon_1^2 + \dots + \varepsilon_d^2} \\ &\geq \frac{1}{\sqrt{d-1}} \|\mathcal{X} - \mathcal{X}_{\text{SVD}}\| \end{aligned}$$

□

EFY. Develop a variant of TT-SVD, which for a prescribed accuracy $\varepsilon > 0$ determines the TT ranks adaptively such that $\|\mathcal{X} - \mathcal{X}_{\text{SVD}}\| \leq \varepsilon$.

Orthogonal TT decompositions

The TT decomposition constructed by TT-SVD satisfies orthogonality relations

$$(U_{\mu}^L)^T U_{\mu}^L = I_{r_{\mu}}, \quad X_{\leq \mu}^T X_{\leq \mu} = I_{r_{\mu}} \quad \text{for } \mu = 1, \dots, d-1.$$

Such a TT decomposition is called **left-orthogonal**.

EFY. Show that $\|\mathcal{X}\| = \|\mathcal{U}_d\|$ holds for a left-orthogonal TT decomposition. What can you say about the singular values of $X^{<\mu-1>}$?

A TT decomposition is called **right-orthogonal** if

$$U_{\mu}^R (U_{\mu}^R)^T = I_{r_{\mu}}, \quad X_{\geq \mu} X_{\geq \mu}^T = I_{r_{\mu}} \quad \text{for } \mu = 2, \dots, d.$$

EFY. Show that $\|\mathcal{X}\| = \|\mathcal{U}_1\|$ holds for a right-orthogonal TT decomposition. What can you say about the singular values of $X^{<1>}$?

Benefits of orthogonal decompositions

- ▶ Numerical stability (see, e.g., Section 3.7 of [Higham'2002] on stability of matrix products).
- ▶ Acceleration of operations.

Orthogonal TT decompositions

Given TT decomposition with cores $\mathcal{U}_1, \dots, \mathcal{U}_d$, can always compute left-/right-orthogonal TT decomposition for the same tensor.

Algorithm to obtain left-orthogonal TT decomposition:

```
for  $\mu = 1, 2, \dots, d - 1$  do  
  Compute QR decomposition  $U_\mu^L = QR$ .  
  Set  $U_\mu^L \leftarrow Q$ .  
  Update  $U_{\mu+1} \leftarrow R \circ_1 U_{\mu+1}$ .  
end for
```

Algorithm to obtain right-orthogonal TT decomposition:

```
for  $\mu = d, d - 1, \dots, 2$  do  
  Compute QR decomposition  $(U_\mu^R)^T = QR$ .  
  Set  $U_\mu^R \leftarrow Q^T$ .  
  Update  $U_{\mu-1} \leftarrow R \circ_3 U_{\mu-1}$ .  
end for
```


Recompression of TT decomposition

Task: Approximate TT decomposition with TT cores $\mathcal{U}_1, \dots, \mathcal{U}_d$ having relatively large TT ranks (R_1, \dots, R_d) by TT decomposition with TT cores $\mathcal{V}_1, \dots, \mathcal{V}_d$ having smaller TT ranks (r_1, \dots, r_d) .

TT-SVD starts with SVD of

$$X^{<1>} = U_1^L X_{\geq 2}^T.$$

Observation: If \mathcal{X} is right-orthogonal then $X_{\geq 2}$ has orthonormal columns and only need to compute SVD of U_1^L !

Compute rank- r_1 approximation (e.g., via truncated SVD):

$$U_1^L \approx V_1^L \Sigma W^T$$

This defines the first TT core \mathcal{V}_1 of the compressed tensor. Move rest to second core:

$$\mathcal{U}_2 \leftarrow \Sigma W^T \circ_1 \mathcal{U}_2.$$

$X_{\geq 3}$ remains the same and thus still has orthonormal columns!

Recompression of TT decomposition

Input: $\mathcal{X} \in \mathbb{R}^{n_1 \times \dots \times n_d}$ in TT decomposition with TT cores

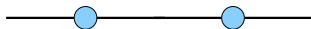
$\mathcal{U}_\mu \in \mathbb{R}^{R_{\mu-1} \times n_\mu \times R_\mu}$, smaller target TT rank (r_1, \dots, r_{d-1}) .

Output: TT cores $\mathcal{V}_\mu \in \mathbb{R}^{r_{\mu-1} \times n_\mu \times r_\mu}$ that define a TT decomposition approximating \mathcal{X} .

- 1: Set $r_0 = r_d = 1$.
 - 2: Right-orthogonalize TT cores \mathcal{U}_μ .
 - 3: **for** $\mu = 1, \dots, d - 1$ **do**
 - 4: Compute rank- r_μ approximation $U_\mu^L \approx V \Sigma W^T$ (e.g., via SVD)
 - 5: Reshape V into $\mathcal{V}_\mu \in \mathbb{R}^{r_{\mu-1} \times n_\mu \times r_\mu}$.
 - 6: Update $\mathcal{U}_{\mu+1} \leftarrow \Sigma W^T \circ_1 \mathcal{U}_{\mu+1}$.
 - 7: **end for**
 - 8: Set $\mathcal{V}_d = \mathcal{U}_d$.
- ▶ Complexity: $\mathcal{O}(dnR^3)$.
 - ▶ Mathematical equivalence to full TT-SVD \rightsquigarrow error bounds carry over.

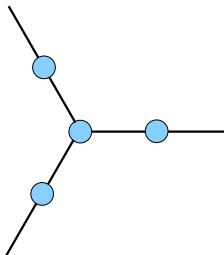
Tensor network diagrams

- ▶ Introduced by Roger Penrose.
- ▶ Heavily used in quantum mechanics (spin networks).
- ▶ Useful to gain intuition and guide design of algorithms.
- ▶ This is the matrix product $C = AB$:



$$C_{ij} = \sum_{k=1}^r A_{ik} B_{kj}$$

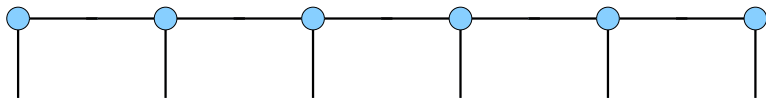
Tensor of order 3 in Tucker decomposition



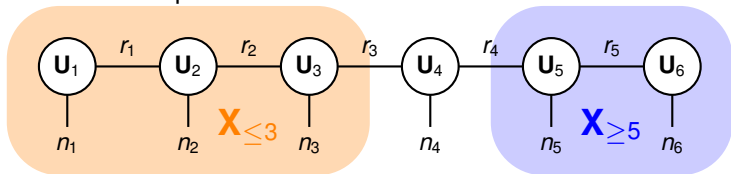
$$\mathcal{X}_{ijk} = \sum_{\ell_1=1}^{r_1} \sum_{\ell_2=1}^{r_2} \sum_{\ell_3=1}^{r_3} \mathcal{C}_{\ell_1 \ell_2 \ell_3} U_{i\ell_1} V_{j\ell_2} W_{k\ell_3}$$

- ▶ $r_1 \times r_2 \times r_3$ core tensor \mathcal{C}
- ▶ $n_1 \times r_1$ matrix U spans first mode
- ▶ $n_2 \times r_2$ matrix V spans second mode
- ▶ $n_3 \times r_3$ matrix W spans third mode.

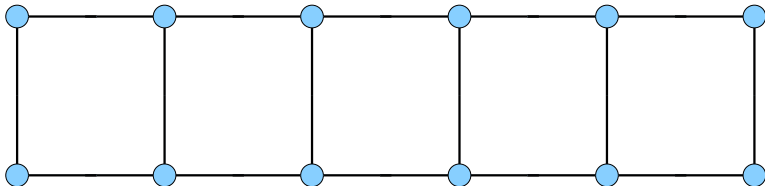
Tensor of order 6 in TT decomposition



- ▶ \mathcal{X} implicitly represented by four $r \times n \times r$ tensors and two $n \times r$ matrices
- ▶ More detailed picture:



Inner product of two tensors in TT decomposition



- ▶ Carrying out contractions requires $O(dnr^4)$ instead of $O(n^d)$ operations for tensors of order d .

EFY. Develop an efficient method for computing the mean,

$$\bar{x} = \frac{1}{n_1 \cdots n_d} \sum_{i_1, \dots, i_d} x(i_1, \dots, i_d),$$

of a tensor in TT decomposition.

Linear operators in TT decomposition

Consider linear operator

$$\mathcal{L} : \mathbb{R}^{n_1 \times \dots \times n_d} \rightarrow \mathbb{R}^{n_1 \times \dots \times n_d}$$

and corresponding matrix representation $\mathcal{L}((i_1, \dots, i_d), (j_1, \dots, j_d))$.

\mathcal{L} is TT operator if it takes the form

$$\begin{aligned} \mathcal{L}((i_1, \dots, i_d), (j_1, \dots, j_d)) = \\ \sum_{k_1=1}^{R_1} \dots \sum_{k_{d-1}=1}^{R_{d-1}} \mathcal{A}_1(1, i_1, j_1, k_1) \mathcal{A}_2(k_1, i_2, j_2, k_2) \dots \mathcal{A}_d(k_{d-1}, i_d, j_d, 1) \end{aligned}$$

with $R_{\mu-1} \times n_{\mu} \times n_{\mu} \times R_{\mu}$ tensors \mathcal{A}_{μ} .

Operator TT rank of \mathcal{L} = Smallest possible values of $(R_1, \dots, R_{\mu-1})$.

Matrix Product Operator (MPO):

$$\mathcal{L}((i_1, \dots, i_d), (j_1, \dots, j_d)) = \mathcal{A}_1(i_1, j_1) \mathcal{A}_2(i_2, j_2) \dots \mathcal{A}_d(i_d, j_d),$$

with $\mathcal{A}_{\mu}(i_{\mu}, j_{\mu}) = \mathcal{A}_{\mu}(:, i_{\mu}, j_{\mu}, :)$.

Linear operators in TT decomposition

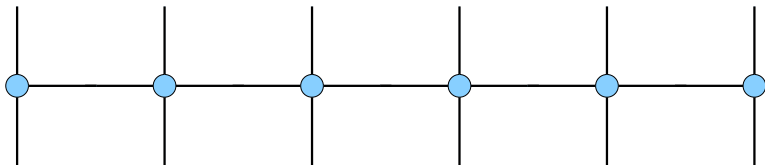
Alternative view: Reinterpret matrix representation of \mathcal{A} as tensor

$$\text{ten}(\mathcal{A}) \in \mathbb{R}^{n_1^2 \times \dots \times n_d^2},$$

by merging indices (i_μ, j_μ) applying TT format, and separating indices (i_μ, j_μ) .

In practice: Perfect shuffle permutation of modes.

Tensor network diagram for operator TT decomposition:



EFY. Show that the identity matrix has operator TT ranks 1.

Linear operators in TT decomposition

Example: Discrete Laplace-like operator takes the form

$$\mathcal{L} : \mathcal{X} \mapsto L_1 \circ_1 \mathcal{X} + \cdots + L_d \circ_d \mathcal{X}$$

for matrices $L_\mu \in \mathbb{R}^{n_\mu \times n_\mu}$. Matrix representation

$$\mathcal{L} = I_{n_d} \otimes \cdots \otimes I_{n_2} \otimes L_1 + \cdots + L_d \otimes I_{n_{d-1}} \otimes \cdots \otimes I_{n_1}$$

TT representation with cores

$$\begin{aligned} A_1(i_1, j_1) &= \begin{pmatrix} L_1(i_1, j_1) & I_{n_1}(i_1, j_1) \end{pmatrix} \\ A_\mu(i_\mu, j_\mu) &= \begin{pmatrix} I_{n_\mu}(i_\mu, j_\mu) & 0 \\ L_\mu(i_\mu, j_\mu) & I_{n_\mu}(i_\mu, j_\mu) \end{pmatrix} \\ A_d(i_d, j_d) &= \begin{pmatrix} I_{n_d}(i_d, j_d) \\ L_d(i_d, j_d) \end{pmatrix} \end{aligned}$$

Multiplication with linear operators in TT decomposition

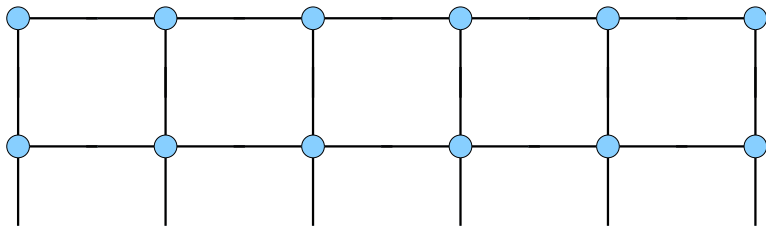
Consider matrix-vector product $\tilde{x} = \mathcal{L}x$ or, equivalently, $\tilde{\mathcal{X}} = \mathcal{L}(\mathcal{X})$.
Assume that \mathcal{L} and \mathcal{U} are in TT decomposition \rightsquigarrow

$$\begin{aligned} & \tilde{\mathcal{X}}(i_1, \dots, i_d) \\ = & \sum_{j_1, \dots, j_d} \mathcal{L}((i_1, \dots, i_d), (j_1, \dots, j_d)) \mathcal{X}(j_1, \dots, j_d) \\ = & \sum_{j_1, \dots, j_d} A_1(i_1, j_1) \cdots A_d(i_d, j_d) U_1(j_1) \cdots U_d(j_d) \\ = & \sum_{j_1, \dots, j_d} (A_1(i_1, j_1) \otimes U_1(j_1)) \cdots (A_d(i_d, j_d) \otimes U_d(j_d)) \\ = & \left[\sum_{j_1} A_1(i_1, j_1) \otimes U_1(j_1) \right] \cdots \left[\sum_{j_d} A_d(i_d, j_d) \otimes U_d(j_d) \right] \\ =: & \tilde{U}_1(i_1) \cdots \tilde{U}_d(i_d), \end{aligned}$$

that is, TT ranks of $\tilde{\mathcal{X}}$ are bounded by TT ranks of \mathcal{X} multiplied with TT operator ranks of \mathcal{L} .

Multiplication with linear operators in TT decomposition

Illustration by tensor network diagrams:



- ▶ Carrying out contractions requires $O(dnr^4)$ operations for tensors of order d .

TT decomposition – Summary of operations

Easy:

- ▶ (partial) contractions
- ▶ multiplication with TT operators
- ▶ recompression/truncation

Medium:

- ▶ entrywise products

Hard:

- ▶ almost everything else

Software:

- ▶ TT toolbox (Matlab, Python), TTeMPS (Matlab), ...